

# Modular Arithmetic

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# 1 Basics

$$a \equiv b \pmod{m} \leftrightarrow a - b = \lambda m, \lambda \in \mathbb{Z}$$

1. Reflexivity
2. Symmetry
3. Transitivity

if  $a_1 \equiv b_1 \pmod{m}, a_2 \equiv b_2 \pmod{m}, a_3 \equiv b_3 \pmod{m}$  then:

- $a + k \equiv b + k \pmod{m}$
- $ak \equiv bk \pmod{m}$
- $a_1 + a_2 \equiv b_1 + b_2 \pmod{m}$
- $a_1 a_2 \equiv b_1 b_2 \pmod{m}$
- $a^k \equiv b^k \pmod{m}$
- $a + k \equiv b + k \pmod{m} \rightarrow a \equiv b \pmod{m}$
- $ak \equiv bk \pmod{m} \wedge \gcd(k, m) = 1 \rightarrow a \equiv b \pmod{m}$
- $a^{-1} \leftrightarrow \gcd(a, m) = 1$
- $a^{-1} \equiv b^{-1} \pmod{m}$
- $ax \equiv b \pmod{m} \rightarrow x \equiv a^{-1}b \pmod{m}$
- $a^{\phi(m)} \equiv 1 \pmod{m}$
- $(p-1)! \equiv -1 \pmod{m}$

## 2 Great Common Divisor (GCD)

Euclides Algorithm is used to solve this problem.

$$a = bq + r$$

$$\gcd(a, b) = \gcd(b, r) = \dots = \gcd(c, 1) = c$$

## 3 Least Common Multiple (LCM)

We can calculate the LCM using GCD:

$$\text{lcm}(a, b) = ab / \gcd(a, b)$$

## 4 Bezout's Theorem

Using Euclides Algorithm we are able to find  $u$  and  $v$ , where:

$$\gcd(a, b) = au + bv : u, v \in \mathbb{Z}$$

## 5 Number Decomposition

This algorithm could be optimized because we only need to look numbers below  $\sqrt{n}$ , and thus, the time is reduced.

## 6 $a^b \pmod{m}$

we are not always to compute high numbers with big exponents, and hence we need another way that simplifies:

we write  $b$  in binary, such as:  $\alpha_1\alpha_2 \dots \alpha_n$ , where  $\alpha_i$  is 0 or 1.

$$a^b \pmod{m} = a^{\alpha_1} \dots a^{\alpha_n} \pmod{m} \text{ and } a^{\alpha_{n+1}} \pmod{m} = (a^{\alpha_n})^2 \pmod{m}$$

## 7 Primality Test

This algorithm tells you, whether a number is prime or not (optimized).

## 8 Find next Prime

This function is important in cryptography. It uses an initial number and finds the next prime over the number given.

## 9 Euler's Totient Function

We need the number decomposition algorithm mentioned before.

$$\phi(n) = (p_1 - 1)p_1^{t_1 - 1} * \dots * (p_n - 1)p_n^{t_n - 1}$$

## 10 Linear Equation Solver

$$ax \equiv b \pmod{m}, b | \gcd(a, m)$$

To solve the equation, we divide by  $d := \gcd(a, m)$

$$a'dx \equiv b'd \pmod{m'd} \rightarrow a'x \equiv b' \pmod{m'}$$

We use Bezout:

$$a'u + m'v = 1$$

Then, we multiply by  $b'$ :

$$b' = a'bu + m'vb \equiv a'du \pmod{m}$$

$$x = du$$

$$x_0 = du + 0m', x_1 = du + 1m', \dots, x_{d-1} = du + (d-1)m'$$

There are  $d$  solutions for this equation

## 11 Inverse

The inverse of a number in mod  $m$  is calculated using the linear solver:

$$a^{-1}x \equiv 1 \pmod{m}$$

## 12 System of Equations

Solving a system of equations is done by constantly solving two equations.

$$\begin{cases} a_1x \equiv b_1 \pmod{m_1} \\ a_2x \equiv b_2 \pmod{m_2} \end{cases}$$

$$\gcd(m_1, m_2) = 1 = m_1u + m_2v$$

$x = m_1ux_2 + m_2vx_1$ , where  $x_1$  and  $x_2$  are the solutions for each equation.

We repeat the process:

$$\begin{cases} x \equiv b \pmod{m_1m_2} \\ a_3 \equiv b_3 \pmod{m_3} \end{cases}$$